



# NTRU+ 규격과 표준화 방안

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## ❖ KeyGen( $1^\lambda$ )

- ◆  $(pk, sk) = \mathbf{Gen}(1^\lambda)$ 
  - $\mathbf{f}', \mathbf{g} \leftarrow \psi_1^n$
  - $\mathbf{f} = 3\mathbf{f}' + \mathbf{1}$
  - check if  $\mathbf{f}$  and  $\mathbf{g}$  are invertible
  - $(pk, sk) = (3\mathbf{g}\mathbf{f}^{-1}, \mathbf{f})$

- SOTP using CBD sampling
- ACWC<sub>2</sub> transform
- FO without re-encryption
- Explicit rejection  $\perp$

## ❖ Encap

- ◆  $m \leftarrow \{0,1\}^n$
- ◆  $(K, \mathbf{r}) = \mathbf{H}(m)$
- ◆  $\mathbf{c} = \mathbf{Enc}'(pk, m; \mathbf{r})$ 
  - $\mathbf{m} = \mathbf{SOTP}(m, \mathbf{G}(\mathbf{r}))$
  - $\mathbf{c} = \mathbf{Enc}(\mathbf{h}, \mathbf{m}; \mathbf{r})$

## ❖ Decap

- ◆  $\mathbf{m}' = \mathbf{Dec}'(sk, \mathbf{c})$ 
  - $\mathbf{m}' = \mathbf{Dec}(\mathbf{f}, \mathbf{c})$
  - $\mathbf{r}' = \mathbf{RRec}(\mathbf{h}, \mathbf{c}, \mathbf{m}')$
  - $m' = \mathbf{Inv}(\mathbf{m}', \mathbf{G}(\mathbf{r}'))$
- ◆  $(K', \mathbf{r}'') = \mathbf{H}(m')$
- ◆ If  $\mathbf{r}' == \mathbf{r}''$ 
  - Return  $K'$
  - Else, return  $\perp$

### ❖ FO(Fujisaki-Okamoto) transform for KEM

- OW/CPA-secure PKE  $\rightarrow$  CCA-secure KEM

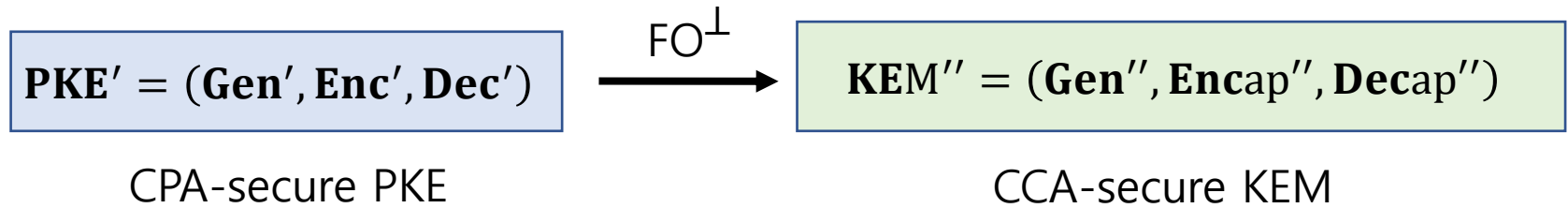
### ❖ Requirements for underlying OW/CPA-secure PKE

- Worst-case correctness error
- Message space with at least 256 min-entropy
- Depending on whether  $\gamma$ -spreadness holds,
  - Two types of decapsulation algorithms exist  $\rightarrow \text{FO}^\perp, \text{FO}^\downarrow$

- (Statistical)  **$\gamma$ -spreadness**

- For fixed message  $m$  and ciphertext  $c$ ,  $\Pr_{r \leftarrow \psi_{\mathcal{R}}} [c = \text{Enc}(pk, m; r)] \leq 2^{-\gamma}$

## ❖ CPA-secure PKE meets $\gamma$ -spreadness ( $\rightarrow$ NTRU+)



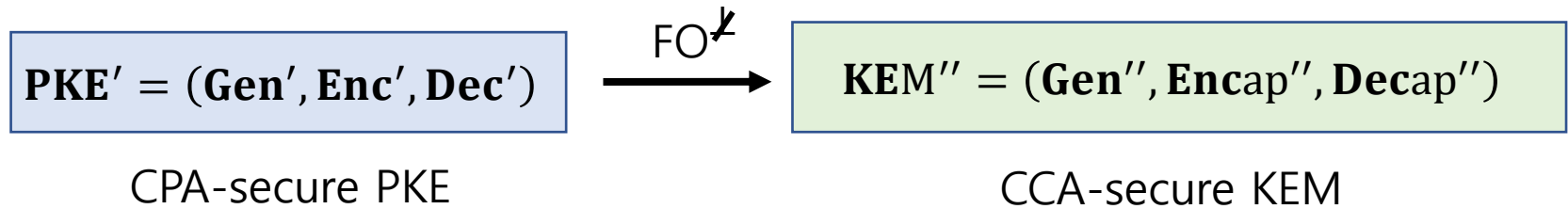
### ◆ $\text{Encap}''(pk)$

- $m \leftarrow \{0,1\}^n$
- $(K, r) = \mathbf{H}(m)$
- $c = \mathbf{Enc}'(pk, m; r)$
- **return**  $(c, K)$

### ◆ $\text{Decap}''(sk, c)$

- $m' = \mathbf{Dec}'(sk, c)$
  - $(K', r') = \mathbf{H}(m')$
  - $c = \mathbf{Enc}'(pk, m'; r')$
  - If  $c = c'$ , return  $K'$ . Else, return  $\perp$
- An orange box labeled "Re-encryption" has a blue arrow pointing to the step  $c = \mathbf{Enc}'(pk, m'; r')$ .

- ❖ CPA-secure PKE does not meet  $\gamma$ -spreadness ( $\rightarrow$  Kyber, SMAUG-T)



◆  $\text{Encap}''(pk)$

- $m \leftarrow \{0,1\}^n$
- $(K, r) = \mathbf{H}(m)$
- $c = \mathbf{Enc}'(pk, m; r)$
- **return**  $(c, K)$

◆  $\text{Decap}''(sk, c)$

- $m' = \mathbf{Dec}'(sk, c)$
- $(K', r') = \mathbf{H}(m')$
- $c = \mathbf{Enc}'(pk, m'; r')$
- $K'' = \mathbf{PRF}(k, c)$
- If  $c = c'$ , return  $K'$ . Else, return  $K''$

Re-encryption

## ❖ Ring Structure

♦  $R_q = \mathbb{Z}_q[x]/\langle \Phi_{3n}(x) \rangle$

▪  $\Phi_{3n}(x) = x^n - x^{n/2} + 1$  :  $3n$ -th cyclotomic polynomial

•  $n = 2^i 3^j$

»  $n = 512, \mathbf{576}, 648, \mathbf{768}, \mathbf{864}, 972, 1024, \mathbf{1152}, \dots$

	Sec. level	n	q	PK (Byte)	CT (Byte)	SK (Byte)	Dec. Failure	Classical (Core-SVP)	Quantum
NTRU+576	1	576	3457	864	864	1,760	$2^{-487}$	111	99
NTRU+768	1+	768		1,152	1,152	2,336	$2^{-379}$	156	139
NTRU+864	3	864		1,296	1,296	2,624	$2^{-340}$	179	160
NTRU+1152	5	1152		1,728	1,728	3,488	$2^{-260}$	248	222

# Performance Comparison - KEMs

Algorithm	sec. (c)	n	q	PK (Byte)	CT (Byte)	SK (Byte)	$\log_2 \delta$	Optimized C (K Cycles)			AVX2 (K Cycles)		
								Gen	Encap	Decap	Gen	Encap	Decap
NTRU+576	111	576	3,457	864	864	1,760	-487	102	60	58	26	26	16
NTRU+768	156	768		1,152	1,152	2,336	-379	124	80	75	28	32	19
NTRU+864	179	864		1,296	1,296	2,624	-340	147	91	87	30	36	22
NTRU+1152	248	1,152		1,728	1,728	3,488	-260	226	118	119	46	47	29
Kyber512	115	256x2	3,329	800	768	1,632	-139	132	157	199	27	29	30
Kyber768	174	256x3		1,184	1,088	2,400	-164	231	256	316	45	44	46
Kyber1024	241	256x4		1,568	1,568	3,168	-174	337	387	464	62	63	67
TiMER	112	256x2	1,024	672	608	832	-161	117	102	137	42	25	35
SMAUG-T 128	112	256x2	1,024	672	672	832	-118	116	103	136	42	25	34
SMAUG-T 192	174	256x3	2,048	1,088	992	1,312	-179	222	205	254	60	47	61
SMAUG-T 256	236	256x4	2,048	1,440	1,374	1,792	-194	357	337	401	82	66	85

**Test Environment:** Intel Core i7-8700K @ 3.70GHz (Single Core, Hyper-Threading Disabled), 64GB RAM, Ubuntu 22.04 LTS, GCC 11.4

**Lattice Estimator :** <https://github.com/malb/lattice-estimator/tree/5ba00f56dd1086c3a42b98fc596c64907adb96ff>

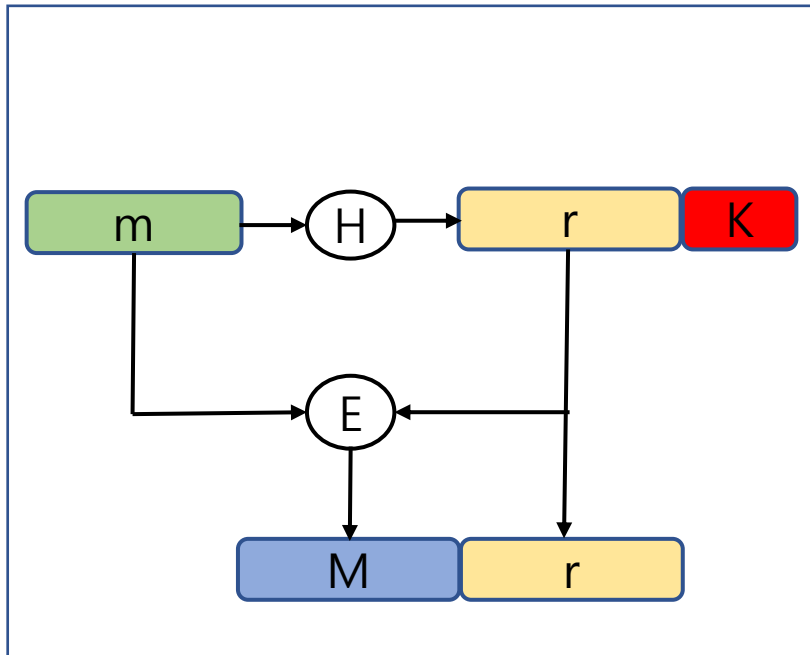
**Kyber source code :** <https://github.com/pq-crystals/kyber/tree/4768bd37c02f9c40a46cb49d4d1f4d5e612bbb882>

**SMAUG-T source code :** [https://github.com/hmchoe0528/SMAUG-T\\_public/tree/1f451c2232cd4e5afd408f2ebacbddb98daef200](https://github.com/hmchoe0528/SMAUG-T_public/tree/1f451c2232cd4e5afd408f2ebacbddb98daef200)

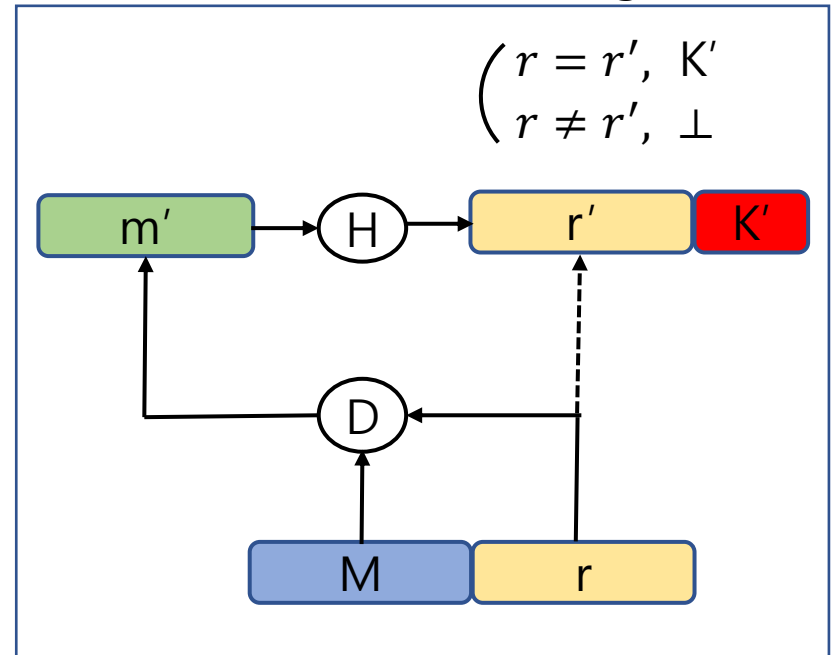
## ❖ OAEP(Optimal Asymmetric Encryption Padding) for KEM

- Using the fact that NTRU+ is **deterministic PKE**

<OAEP encoding>



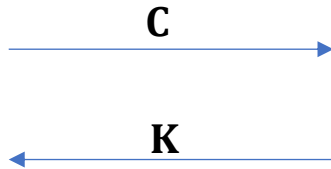
<OAEP decoding>



- $H$ : hash function
- $SOTP = (E, D)$



## ❖ Using decap. Oracles + side-channel analysis



### ◆ $\text{Decap}''(sk, c)$

- $m' = \text{Dec}'(sk, c)$
- $(K', r') = H(m')$
- $c = \text{Enc}'(pk, m'; r')$
- $K'' = \text{PRF}(k, c)$
- If  $c = c'$ , return  $K'$ . Else, return  $K''$

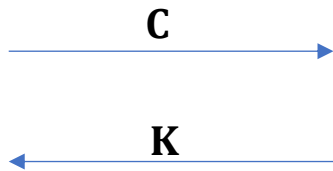
Side-channel info.



## ❖ In chosen-ciphertext attacks on $C = \text{Encap}''(pk, m)$

- Side-channel info. = same  $m$  ???  $\rightarrow$  private key( $sk$ ) recovery !

## ❖ Using decap. Oracles + side-channel analysis



### ◆ Decap''( $sk, c$ )

- $m' = \text{Dec}'(sk, c)$

- $m' = \text{Dec}(\mathbf{f}, c)$

- $r' = \text{RRec}(h, c, m')$

- $m' = \text{Inv}(m', G(r'))$

- $(K', r'') = H(m')$

- If  $r' == r''$

- Return  $K'$ , else return  $\perp$

Side-channel info.



## ❖ In chosen-ciphertext attacks on $C = \text{Encap}''(pk, m)$

- Side-channel info. = same  $r$  ???  $\rightarrow$  message( $m$ ) recovery !

## ❖ Rabin( $sk$ recovery) vs. RSA( $m$ recovery)

# Comparison to Kyber and SMAUG-T

Scheme	SMAUG-T [CCH+23]	Kyber [BDK+18]	NTRU+[KP23]
NTT-friendly	No	Yes	Yes
Correctness error	Worst-case	Worst-case	Worst-case
Message set	$m \leftarrow \{0,1\}^{256}$	$m \leftarrow \{0,1\}^{256}$	$m \leftarrow \{0,1\}^n$
Message distribution	Uniform	Uniform	Arbitrary
CCA transform	$\text{FO}^\nless$	$\text{FO}^\nless$	$\text{ACWC}_2 + \overline{\text{FO}}^\perp$
Hardness assumptions	Module RLWE	Module RLWE	NTRU, RLWE
Re-encryption in Decap.	Yes	Yes	No
In response to invalid CT	PRF key	PRF key	$\perp$
CCA attack + SCA	Private key recovery	Private key recovery	Message recovery

**[CCH+23]** Cheon et al. "SMAUG: Pushing Lattice-Based Key Encapsulation Mechanisms to the Limits", SAC 2023

**[BDK+18]** Bos et al. "CRYSTALS - Kyber: A CCA-Secure Module-Lattice-Based KEM", IEEE EuroS&P 2018

**[KP23]** Kim et al. "NTRU+: Compact Construction of NTRU Using Simple Encoding Method", IEEE TIFS, 18 (2023)

## ❖ KS 표준 문서 작성 일정

	6월	7월	8월	9월	10월	11월
선행 표준화 사례 조사	○————○					
표준 구성 및 작성 계획 수립		○————○				
초안 완성			○————○			
최종안 완성						○————○

## ❖ KS 표준 문서 작성시 고려사항

- ◆ 개념의 명확성을 고려한 용어 및 명칭 사용
- ◆ 국제 표준을 목표로 ISO 등 국제 표준 문서를 참고하여 작성
- ◆ 기존 국제 표준을 참조하여 기법 내의 함수 구성을 정의
- ◆ 비전문가가 이해할 수 있도록 각 세부 알고리즘을 함수로 표현

**T**hank You

**Q&A**